

(u, v) , weight $w(u, v)$

$$w(P) = \sum_{(u, v) \in P} w(u, v)$$

$\delta(u, v) =$ u থেকে v এর মধ্যে শর্টেস্ট পথের লেংথ

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$$w(P) = \delta(u, v) = \infty \text{ যদি } u \not\rightarrow v \quad u \xrightarrow{w(u, v)} v$$

$w(u, v) \in \mathbb{R}$

$w(u, v)$

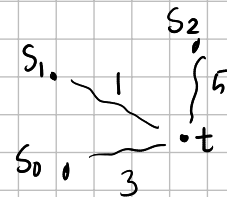
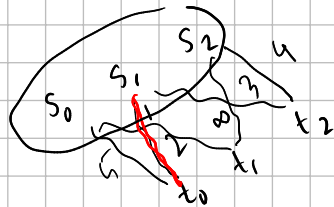
Main Problem: given S (source), find for every other vertex $t \neq s$, $\delta(s, t)$

Variations:

(i) Multi-Source

(ii) Multi-Source \rightarrow Multi destination

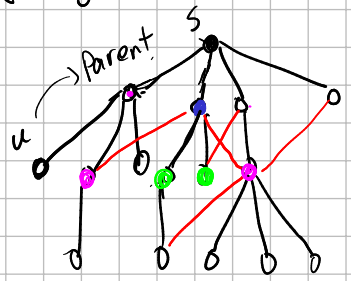
\hookrightarrow weighted
 \hookrightarrow unweighted



$$O(V+E)$$
$$O(E \log V)$$

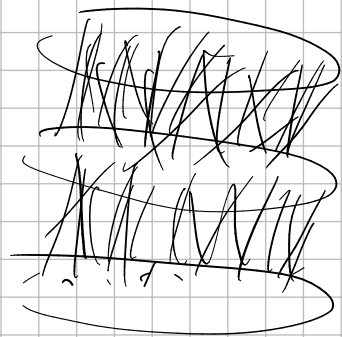
BFS (Single Source)

Level-0
Level-1
Level-2
Level-3



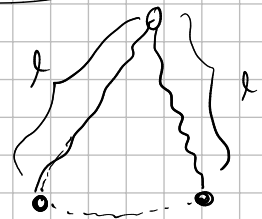
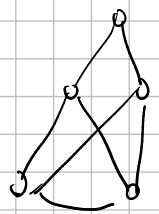
Black = Tree edge
Red = Extra edge

BFS tree



edge in same level \Rightarrow not bipartite
not bipartite \Rightarrow same level \Rightarrow edge
bipartite \Rightarrow no edge \in same level

0
1
2



$2l+1$

BFS (Multisource)

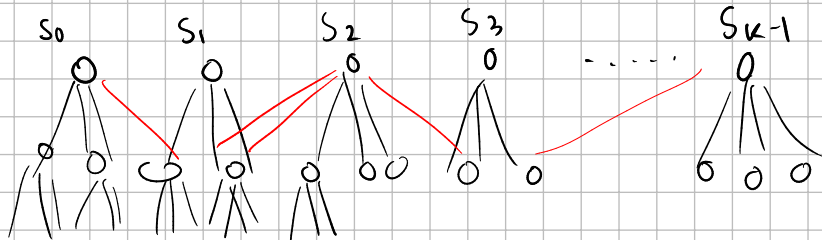
$s_0, s_1, s_2 \dots s_{k-1}$

for each t , calculate $\underline{d[t]} = \min_{x \in S} \{ \underline{\delta(x, t)} \}$

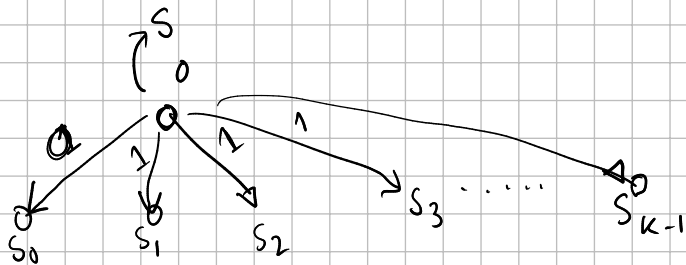
queue

$q \leftarrow s_0, s_1, s_2 \dots s_{k-1}$, $d[s_0] = d[s_1] = d[s_2] = \dots = d[s_{k-1}] = 0$

level



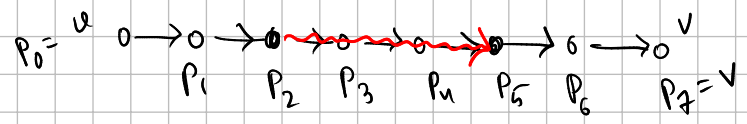
1 2 ... n



Optimal Substructure of a Shortest Path

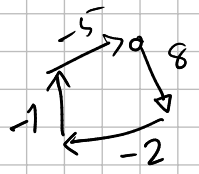
if P is of a shortest path $u \rightarrow v$,

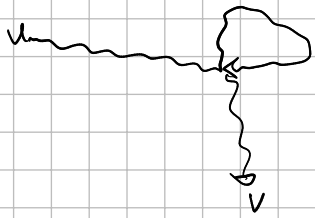
the red path is a shortest path of $P_2 \rightarrow P_5$



Negative weights

Negative Cycle! যেই সাইকেলের weight এর sum negative





$$n = |V|$$

$$\# \text{ edges} \leq |V| - 1$$

$$d[s] = 0$$

$v.p.c$

$$s.p.c = \text{NIL}$$

parent

Relaxation

u, v, w

$$d[\dots] = \infty$$

$$d[s] = 0$$

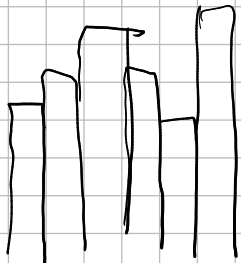
$$d \quad u \xrightarrow{w} v$$

Arbitrarily choose an edge (u, v, w) :
relax (u, v, w)

$$\text{relax}(u, v, w):$$

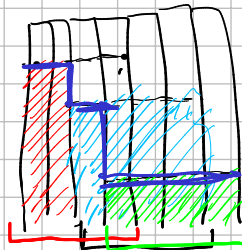
$$d[v] = \min(d[v], d[u] + w)$$

$$d[v] \leq d[u] + w$$



(l, r, x)

$a_i \leq x$ for each i



(l_i, r_i, x_i)

for $(i = l; i \leq r; i++)$
 $a_i = \min(a_i, x)$



$$u \xrightarrow{w} v$$

$d[u]$ = shortest path length from s to u

$$\left. \begin{matrix} 200 \\ \text{max}(a, b) \end{matrix} \right\} d[v] > d[u] + w$$

condition $\left\{ \begin{array}{l} d[v] \leq d[u] + w \\ \text{each edge is not } vq \text{ but a condition} \end{array} \right.$

Single Source Shortest Path (Weighted Graph)

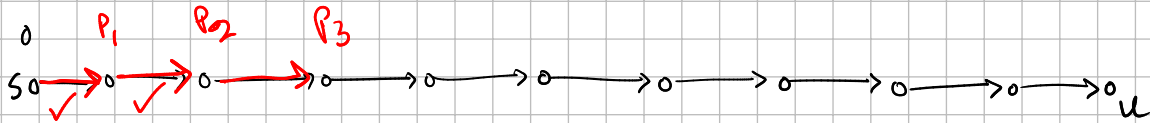
The following algorithm is correct:

- Big $\left[\begin{array}{l} \text{① Iterate on the edges, and find an edge } (u, v, w) \text{ such that } d[v] > d[u] + w \\ \text{and do, } d[v] = d[u] + w \\ \text{② Exit if you cannot find such an edge} \end{array} \right.$

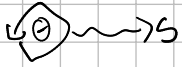
$T \leq |V|$

Spoiler: $T \leq |V| - 1$

$$O(T \times E) \xrightarrow{|V|} O(|V| \times E) \rightarrow \sum_u d[u]$$



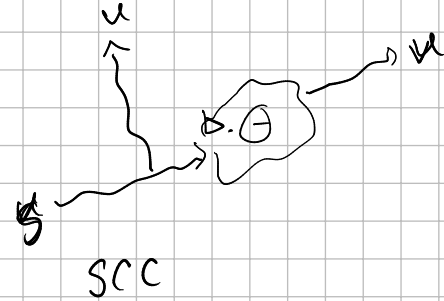
No negative cycle \Rightarrow you cannot find such an edge (u, v, w) that satisfies $d[v] > d[u] + w$ in the (V) -th application of "Big-R".



$$|P| \geq V$$

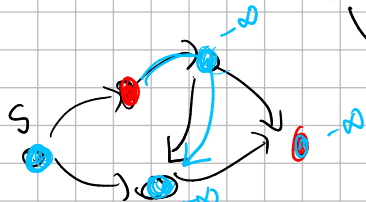
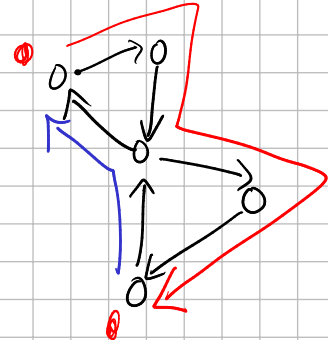
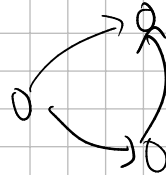
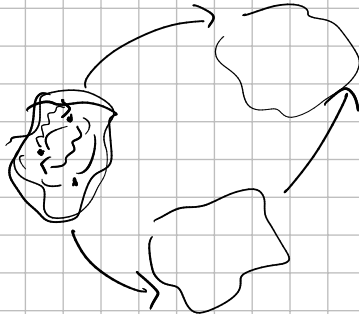


$|V|$

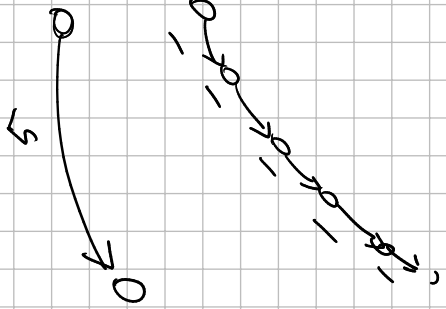


"Source (S) negative cycle reachable"

~~SCC~~



E^2

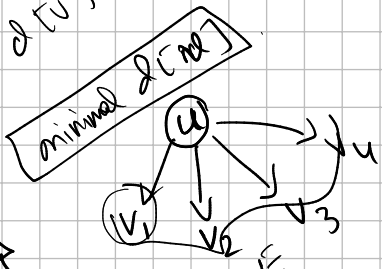
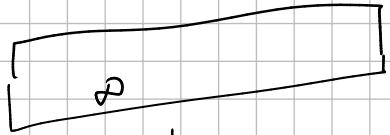
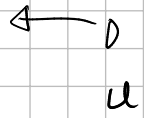


(V)

$$d[v], d[u] + w$$

$$d[v] = d[w] + w$$

$$\min \{ d[u] \}$$



$T (V+E)$

$$(V+E) \lg V$$

- $d[v_1]$
- $d[v_2]$
- $d[v_3]$

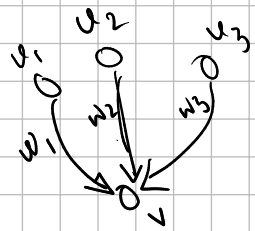
$V^2 + E$

$$\lg + |adj_u| \times \lg$$

$$V \lg$$

$$L, E \lg$$

$$(E/2)^2 \lg$$



$$u_i + w_i = \delta(s, v)$$

DAG