

$(u, v)$ , weight  $w(u, v)$

$$w(P) = \sum_{(u, v) \in P} w(u, v)$$

$\delta(u, v) =$  u থেকে v এর মধ্যে শর্টেস্ট পাতের লেংথ

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$$w(P) = \delta(u, v) = \infty \text{ যদি } u \not\rightarrow v \quad u \xrightarrow{w(u, v)} v$$

$w(u, v) \in \mathbb{R}$

$w(u, v)$

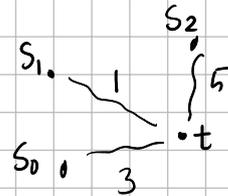
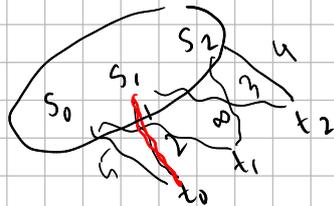
Main Problem: given  $S$  (source), find for every other vertex  $t \neq s$ ,  $\delta(s, t)$

Variations:

(i) Multi-Source

(ii) Multi-Source  $\rightarrow$  Multi destination

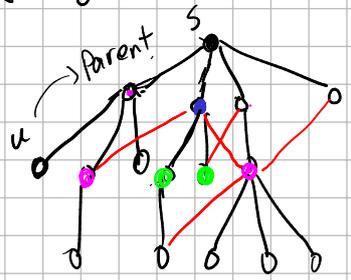
$\hookrightarrow$  weighted  
 $\hookrightarrow$  unweighted



$$O(V+E)$$
$$O(E \log V)$$

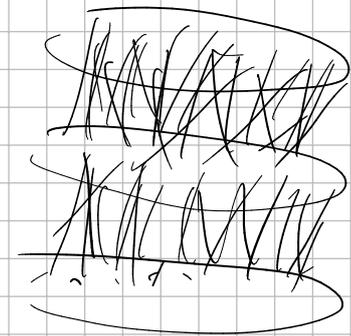
# BFS (Single Source)

Level-0  
Level-1  
Level-2  
Level-3



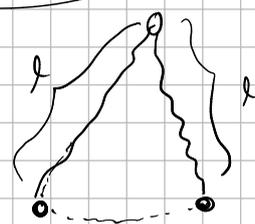
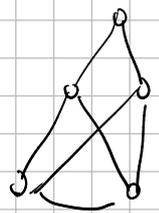
Black = Tree edge  
Red = Extra edge

BFS tree



edge in same level  $\Rightarrow$  not bipartite  
not bipartite  $\Rightarrow$  same level  $\rightarrow$  edge  
bipartite  $\Rightarrow$  no edge  $\in$  same level

0  
1  
2



$2l+1$

# BFS (Multisource)

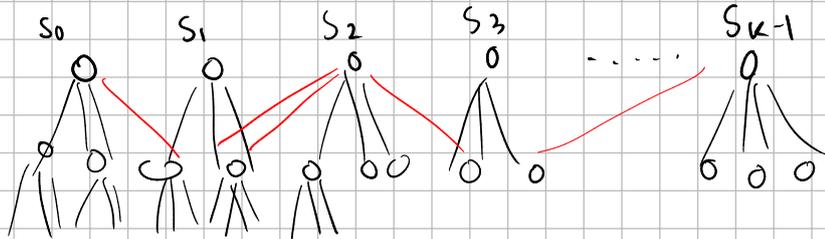
$s_0, s_1, s_2 \dots s_{k-1}$

for each  $t$ , calculate  $\underline{d[t]} = \min_{x \in S} \{ \underline{\delta(x, t)} \}$

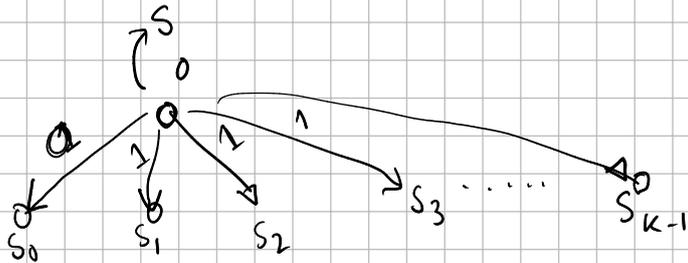
queue

$q \leftarrow s_0, s_1, s_2 \dots s_{k-1}$ ,  $d[s_0] = d[s_1] = d[s_2] = \dots = d[s_{k-1}] = 0$

level



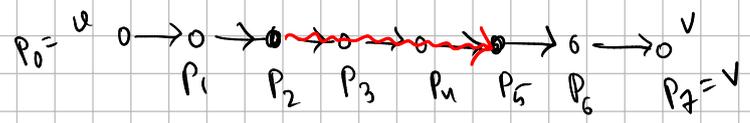
1 2 ... n



# Optimal Substructure of a Shortest Path

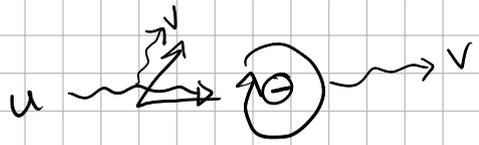
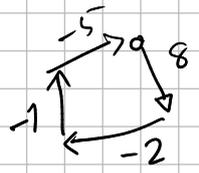
if  $P$  is of a shortest path  $u \rightarrow v$ ,

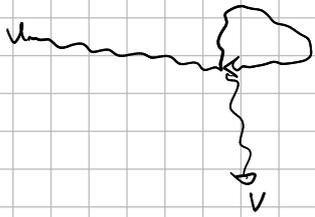
the red path is a shortest path of  $P_2 \rightarrow P_5$



## Negative weights

Negative Cycle! যেই সাইকেলের weight এর sum negative





$$n = |V|$$

$$\# \text{ edges} \leq |V| - 1$$

$$d[s] = 0$$

$v.t.c$

$$s.p.t = NIL$$

parent

Relaxation

$u, v, w$

$$d[\dots] = \infty$$

$$d[s] = 0$$

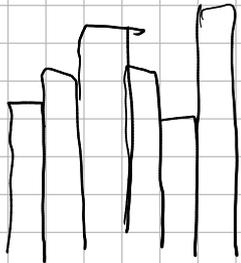
$$d \quad u \xrightarrow{w} v$$

Arbitrarily choose an edge  $(u, v, w)$ :  
relax  $(u, v, w)$

$$\text{relax}(u, v, w):$$

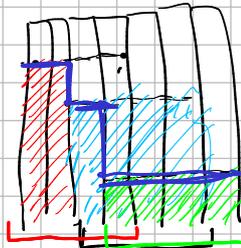
$$d[v] = \min(d[v], d[u] + w)$$

$$d[v] \leq d[u] + w$$



$(l, r, x)$

$a_i \leq x$  for each  $i$



$(l_i, r_i, x_i)$

for  $(i = l; i \leq r; i++)$   
 $a_i = \min(a_i, x)$



$$u \xrightarrow{w} v$$

$d[u]$  = shortest path length from  $s$  to  $u$

$$\left. \begin{matrix} 200 \\ \text{max}(a, b) \end{matrix} \right\} d[v] > d[u] + w$$

condition  $\left\{ \begin{array}{l} d[v] \leq d[u] + w \\ \text{each edge is not } vq \text{ but a condition} \end{array} \right.$

## Single Source Shortest Path (Weighted Graph)

The following algorithm is correct:

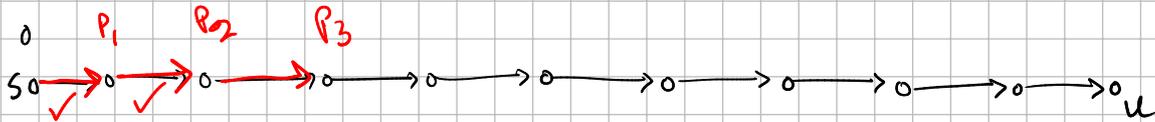
- Big  $\left[ \begin{array}{l} \text{① Iterate on the edges, and find an edge } (u, v, w) \text{ such that } d[v] > d[u] + w \\ \text{and do, } d[v] = d[u] + w \\ \text{② Exit if you cannot find such an edge} \end{array} \right.$

$T \leq |V| - 1$

Spoiler:  $T \leq |V| - 1$

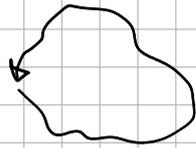
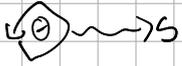
$$O(T \times E) \xrightarrow{\text{red}} O(V \times E) \rightarrow O(V^2)$$

$\left( \sum_u d[u] \right)$



No negative cycle  $\Rightarrow$  you cannot find such an edge  $(u, v, w)$  that satisfies  $d[v] > d[u] + w$  in the  $(V)$ -th application of "Big-R".

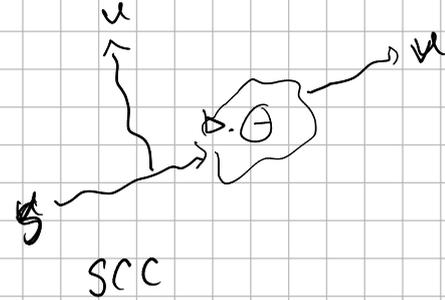
$\Leftarrow$



$|P| \geq V$

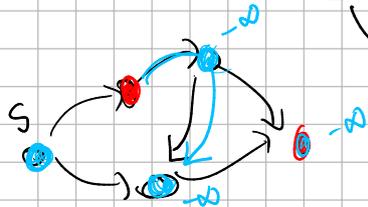
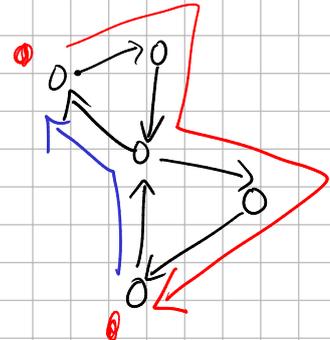
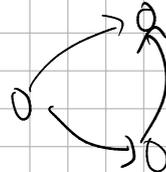
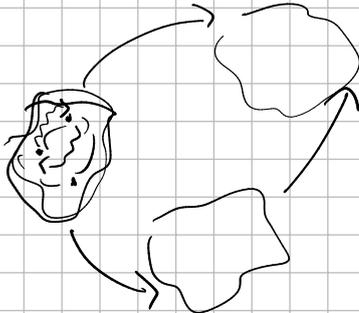


$|V|$

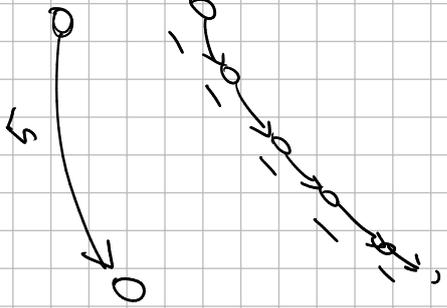


"Source (S) negative cycle reachable"

~~SCC~~



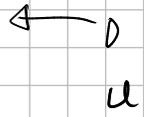
$E^2$



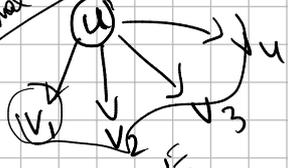
$(V)$

$d[v], d[u] + w$   
 $d[v] = d[w] + w$

$\min \{ d[u] \}$



minimal  $d[u]$



$T (V+E)$

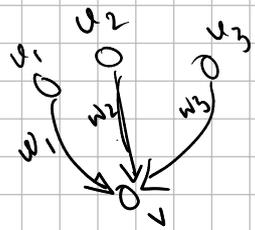
- $d[v_1]$
- $d[v_2]$
- $d[v_3]$

$V+E$

$(V+E) \lg V$   
 $\lg + |adj_u| \times \lg$   
 $V \lg$   
 $L, E \lg$



$(E/2)^2 \lg$



$u_i + w_i = \delta(s, v)$   
 DAG